

# CONTRARIETY VS. POST-COMPLEMENT

Logico-linguistic issues around classical and  
modern squares of opposition

(for Ed Keenan)

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# Outline

- 1 Prologue
- 2 Introduction and Preliminaries
- 3 Square Theory
- 4 Square Examples

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$$(at\ most\ p/q\ of\ the) \neg = at\ least\ (q-p)/q\ of\ the\ (0 < p < q)$$

$$(Q_1 \wedge Q_2) \neg = Q_1 \neg \wedge Q_2 \neg$$
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 and thus, for  $Q = between\ p/q\ and\ (q-p)/q\ of\ the$ ,  $Q_{\neg} = Q$ .
- It illustrates one thing I will argue today: post-complement is, logically as well as linguistically, more interesting than other forms of 'inner' negation proposed, notably **contrariety**.

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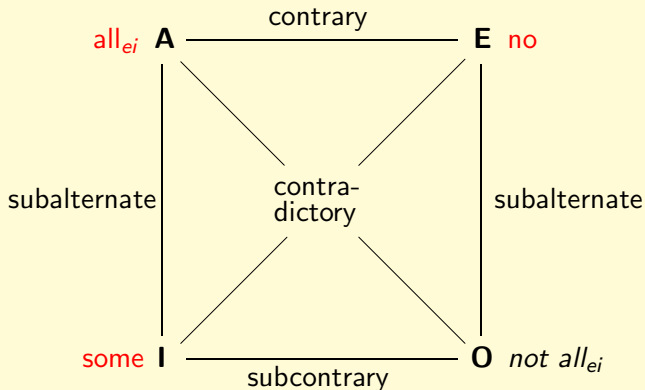
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- Good overviews of this debate are
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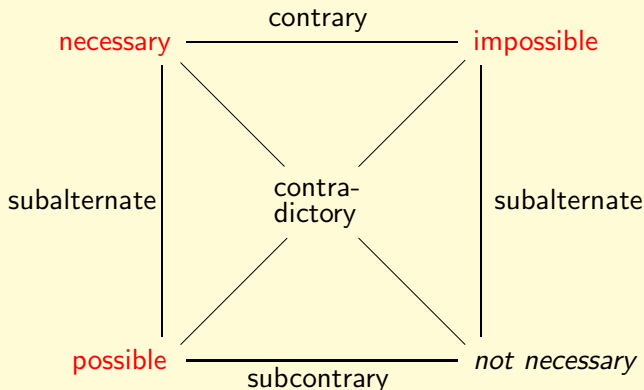
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- Some of the remarks here originate in S. Peters and D. Westerståhl, *Quantifiers in Language and Logic*, OUP, 2006, chs. 1 and 4.

# Aristotle's (quantified) Square



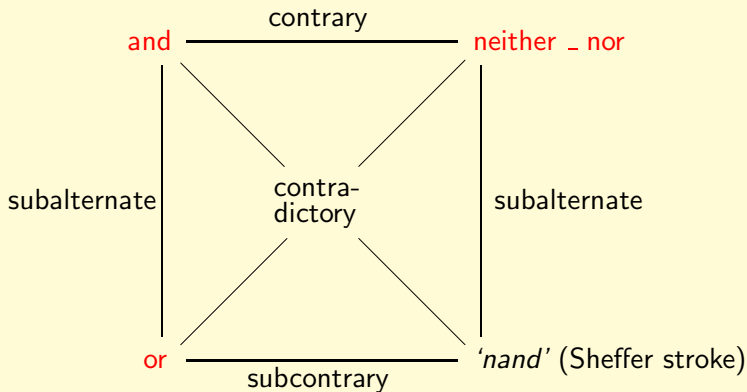
(Red color indicates lexicalization.)

# A modal square (also Aristotle)



And there is a *deontic square* (obligatory, forbidden, permitted, not obligatory), and a *temporal adverb square* (always, never, sometimes, not always), ...

## And a propositional square



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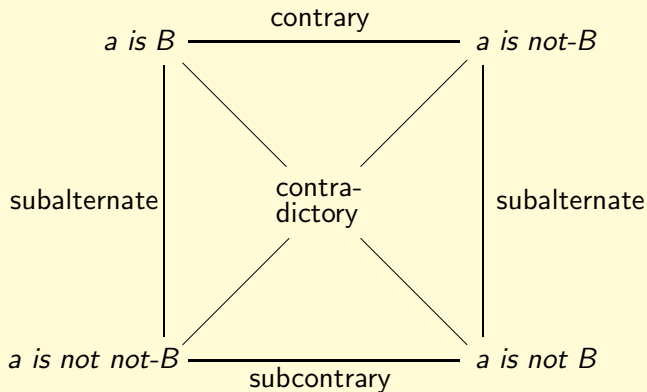
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- However, Aristotle considered yet another square:

# The singular square



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- What happens when 'a' doesn't denote (*Santa Claus is a man*, *The largest prime number is not even*), or when it denotes something outside the range of significance for *B* (*The number 2 is green*)?
- There seem to be different kinds of contraries: cf. *white/not-white* (logical contraries), *white/black* (polar contraries), and *white/red*.



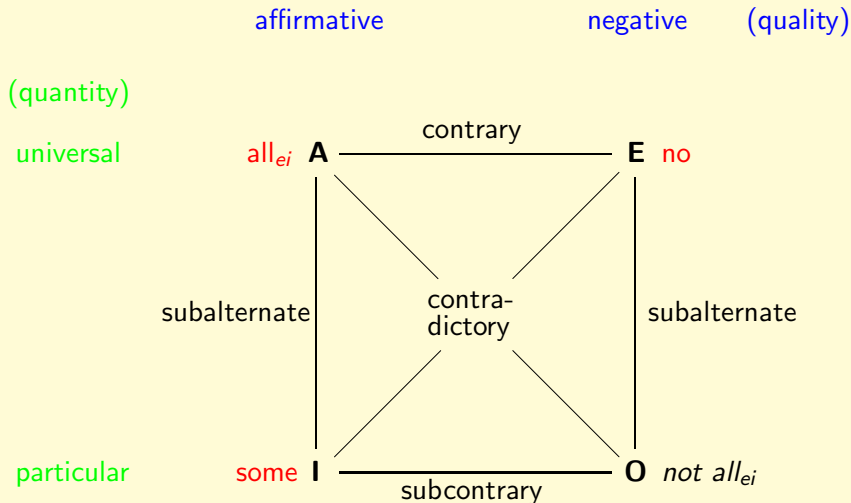
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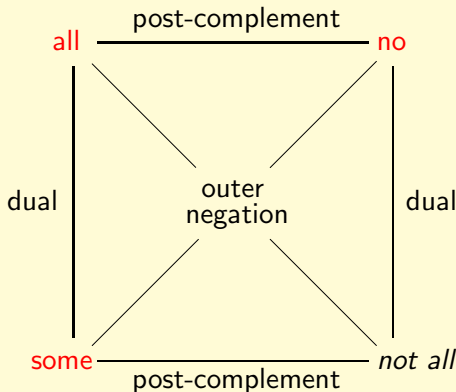
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- So, back to Aristotle's (general, quantified) square:

# Aristotle's Square



# The Modern Square

The **modern square** (spanned by *all*) differs from Aristotle's *only* in the fact that *all* replaces *all<sub>ei</sub>*. But this apparently small difference is in fact huge.



# Logical relations in the squares

- Aristotle's square:

- 1 **contradictoriness** (diagonal) = **outer negation**:  $Q' = \neg Q$
- 2 **contrariety** (upper horizontal):  $Q(A, B) \Rightarrow \neg Q'(A, B)$
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- Modern square:
  - 1 outer negation (diagonal):  $Q' = \neg Q$
  - 2 **inner negation** or **post-complement** (horizontal):  $Q' = Q \neg$   
[where  $Q \neg(A, B) \Leftrightarrow Q(A, M - B) \Leftrightarrow Q(A, \overline{B})$ ]
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- (No quarrel about the meaning of negation here. The quarrel is about which of these relations are fundamental.)
- **Existential import** [ $Q(A, B) \Rightarrow A \neq \emptyset$ ] at *affirmative* positions in Aristotle's square; at *particular* positions in modern square.

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- These issues will not be dealt with here.

## Another main issue:

- I will, however, say something about existential import (EI).  
Three traditional positions (Horn 1989, 1997):
  - 1 EI stems from (affirmative) *quality* (Aristotle (probably), Apuleios, Boethius, Abelard, Carroll, traditional logic).
  - 2 EI stems from (particular) *quantity* (Frege, modern logic).
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- My own position (to be argued) is that (a) for the squares discussed so far, EI is a pragmatic matter (Gricean implicature; this is roughly Horn's position), although the notion of EI is too narrow, but (b) semantic issues of the EI type exist for certain *other squares*.



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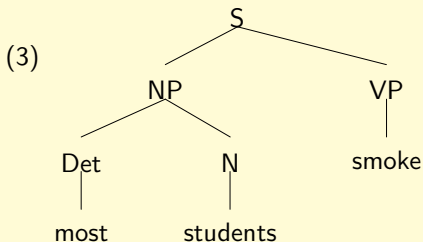
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- So while Aristotle’s (quantificational) square illustrates some of the behavior of a few selected quantifiers, the modern square reveals more general *patterns of quantification*.
- To see this, we need one more preliminary.

## (Generalized) quantifiers

- A (generalized) **quantifier** (of type  $\langle 1, 1 \rangle$ )  $Q$  associates with each universe  $M$  a binary relation  $Q_M$  between subsets of  $M$ .

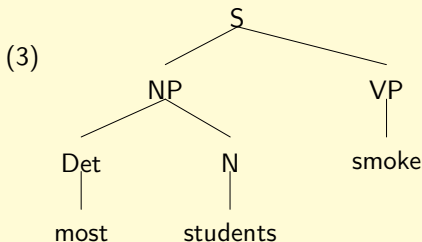
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- Such denotations are CONSERV [ $Q_M(A, B) \Leftrightarrow Q_M(A, A \cap B)$ ] and EXT [if  $A, B \subseteq M \subseteq M'$  then  $Q_M(A, B) \Leftrightarrow Q_{M'}(A, B)$ ], and often ISOM. CONSERV + EXT (as well as ISOM) is *preserved* under the three kinds of negation (but *not* under e.g. contrariness.)



## Examples of quantifiers denoted by determiners:

- (By EXT there is no need to mention  $M$ .)

$$\text{all}(A, B) \iff A \subseteq B$$

$$\text{all}_{\text{ei}}(A, B) \iff \emptyset \neq A \subseteq B$$

$$\text{at least two}(A, B) \iff |A \cap B| \geq 2$$

$$\text{exactly five}(A, B) \iff |A \cap B| = 5$$

$$\text{all but three}(A, B) \iff |A - B| = 3$$

$$\text{more than two thirds of the}(A, B) \iff |A \cap B| > 2/3 \cdot |A|$$

$$\text{most}(A, B) \iff |A \cap B| > |A - B|$$

$$\text{the ten}(A, B) \iff |A| = 10 \text{ and } A \subseteq B$$

$$\text{John's}(A, B) \iff \emptyset \neq A \cap \{a : \text{John 'possesses' } a\} \subseteq B$$

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- Let **1** (**0**) be the trivially true (false) quantifier (*at least zero*, *fewer than zero*).

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- **FACT 1:**
  - (a)  $square(\mathbf{0}) = square(\mathbf{1}) = \{\mathbf{0}, \mathbf{1}\}$ .
  - (b) If  $Q$  is non-trivial, so are the other quantifiers in its square.
  - (c) If  $Q' \in square(Q)$ , then  $square(Q) = square(Q')$ . So any two squares are either identical or disjoint.
  - (d)  $square(Q)$  has either two or four members.

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  - (b) If  $Q$  is non-trivial, so are the other quantifiers in its square.
  - (c) If  $Q' \in square(Q)$ , then  $square(Q) = square(Q')$ . So any two squares are either identical or disjoint.
  - (d)  $square(Q)$  has either two or four members.
- *Ad (d):* Normally four, but consider Keenan's *between one-third and two-thirds of the!*

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- Each position determines the diagonally opposed quantifier, i.e, its outer negation, but *not* the quantifiers at the other two positions. For example:
- **FACT 2:** The square

[**A**: *at least five*; **E**: *no*; **I**: *some*; **O**: *at most four*]  
is classical. More generally, for  $n \geq k$ ,

[**A**: *at least n*; **E**: *fewer than k*; **I**: *at least k*; **O**: *fewer than n*]  
is classical.

# A logico-linguistic project

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- What is the role of existential import (and other kinds of import)?
- Such a project cannot even be formulated with classical squares (since these are not determined).
- Which is not to say that the classical logical oppositions always fail for squares other than  $\{all_{ei}, no, not\ all_{ei}, some\}$ . When do they hold and when don't they?

## Qualitative and quantitative aspects of modern squares?

- Since, for  $Q' \in \text{square}(Q)$ ,  $\text{square}(Q') = \text{square}(Q)$ , have we lost all qualitative and quantitative aspects in modern squares? I.e. can *any* quantifier occupy the **A** corner?

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- Some familiar monotonicity properties:

MON $\uparrow$ :  $Q(A, B) \ \& \ B \subseteq B' \Rightarrow Q(A, B')$

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etc.

- **FACT 3:** The monotonicity behavior of  $Q$  determines the monotonicity behavior of all elements of  $\text{square}(Q)$ :
  - ①  $Q$  is  $\text{MON}\uparrow$  iff  $Q\neg$  is  $\text{MON}\downarrow$  iff  $\neg Q$  is  $\text{MON}\downarrow$  iff  $Q^d$  is  $\text{MON}\uparrow$
  - ②  $Q$  is  $\uparrow\text{MON}$  iff  $Q\neg$  is  $\uparrow\text{MON}$  iff  $\neg Q$  is  $\downarrow\text{MON}$  iff  $Q^d$  is  $\downarrow\text{MON}$
  - ③ So if  $Q$  is *doubly* monotone, all four combinations are exemplified in its square.

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- This would make it possible to identify the exact position in the square of any doubly monotone quantifier.
- However, many quantifiers are only *right monotone*, e.g. the proportional quantifiers. So we would know, for example, that *at least two-thirds of the* is affirmative: either **A** or **I**.
- And this seems right, since the dual of *at least two-thirds of the* is *more than one-third of the*, and it seems arbitrary which of these two should go into the **A** position.

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  - 1  $Q$  is symmetric iff  $Q \neg$  is co-symmetric iff  $\neg Q$  symmetric iff  $Q^d$  is co-symmetric
  - 2 Also, under CONSERV, symmetry is the same as *intersectivity* (Keenan's term): if  $A \cap B = A' \cap B'$  then  $Q(A, B) \Leftrightarrow Q(A', B')$ . So co-symmetry = co-intersectivity: if  $A - B = A' - B'$  then  $Q(A, B) \Leftrightarrow Q(A', B')$ .

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- For example, *at most ten* is at the **E** position.
- But we already knew that, since *at most ten* is  $\downarrow\text{MON}\downarrow$ .  
Indeed, if  $Q$  is right monotone and symmetric, it is clearly also left monotone.
- The cases where symmetry would give extra information are rather limited. This follows from the next result.

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- So we get extra information for cases like *an even number of*, which is logical and symmetric but not right monotone, and *no – except John*, which is non-logical (CONSERV and EXT, but not ISOM) and symmetric but not right monotone.
- I.e. there are two possible configurations of *square(an even number of)* and *square(no – except John)*.

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- **FACT 9:** If  $Q$  is logical and  $\uparrow\text{MON}$ , it is a finite disjunction of quantifiers of the form  $|A \cap B| \geq n \ \& \ |A - B| \geq k$ .



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- On the other hand, *most of John's* is  $\text{MON}\uparrow$  but not left monotone, so it is affirmative, but there seems to be no logical indication of whether it is **A** or **I**.

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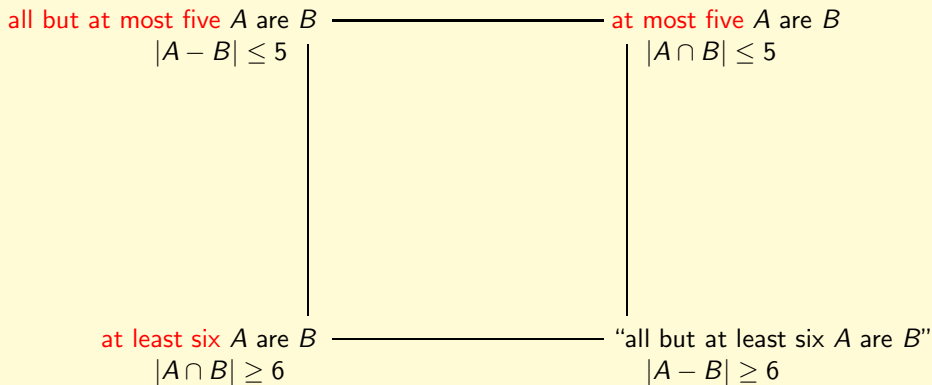
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- With this, it's time to look at a few examples.

## Example I: numerical quantifiers 1

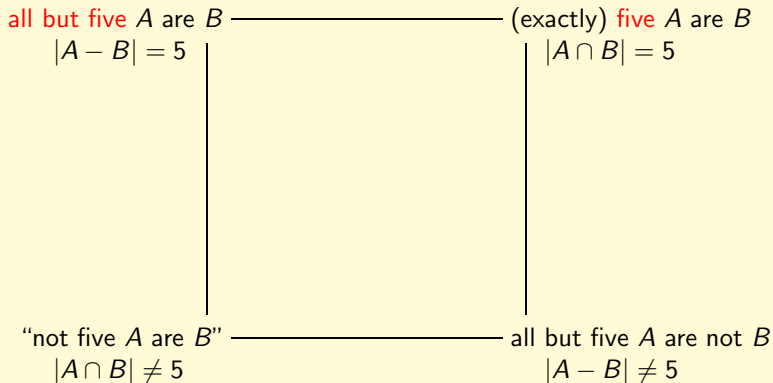
$Q = \text{at least six}, (\uparrow \text{MON} \uparrow)$

(Now red color means 'realizable' as a (possible complex) determiner.)



## Example I: numerical quantifiers 2

$Q = (\text{exactly}) \text{ five}$ . Symmetric, not monotone, so two configurations possible. Choose the one below, since for  $n = 0$  it becomes the (modern) Aristotelian square.



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- But here we would need very odd ‘existential imports’:
- **FACT 10:**
  - (a) *square(at least  $n+1$ )* is classical iff  $|A| > 2n$  is presupposed.
  - (b) *square(exactly  $n$ )* is classical iff  $|A| \neq 2n$  is presupposed.

## Digression: standard EI too weak even for Aristotelian square

- Compare:
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- These are normally fine, but odd to astronomers, who know that Mars has just two moons.
- Conclusion: relegate facts like these to pragmatics.



## Example II: proportional quantifiers

$Q = \text{at least } 2/3 \text{ of the. MON}\uparrow$ , but not left monotone or symmetric, so two configurations possible. No obvious choice between them, i.e. whether  $Q$  or  $Q^d$  is **A**.

at least  $2/3$  of the  $A$  are  $B$

$$|A \cap B| \geq 2/3 \cdot |A|$$

at most  $1/3$  of the  $A$  are  $B$

$$|A \cap B| \leq 1/3 \cdot |A|$$

more than  $1/3$  of the  $A$  are  $B$

$$|A \cap B| > 1/3 \cdot |A|$$

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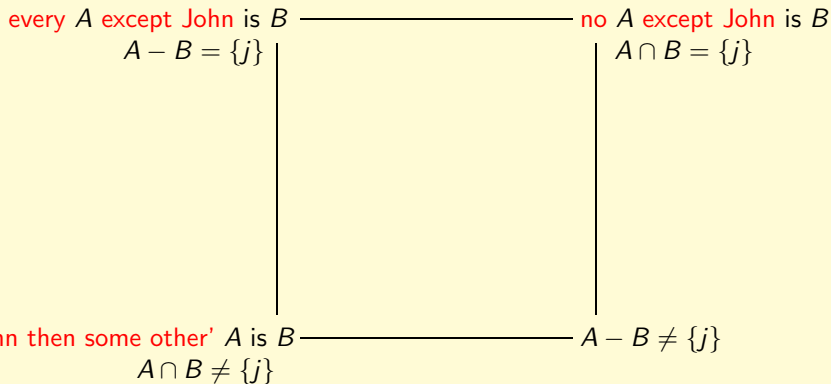
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- If we do, this square is both classical and modern.

## Example III: an 'exceptive' square

$Q = no \text{ } \_ \text{ except John}$ . Symmetric, not monotone; two configurations possible. Is  $Q$  **I** or **E**? The latter choice makes the **O** corner 'unrealized' (as usual), and reduces to the (modern) Aristotelian square when the exception set is empty:





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  - 2 'Freedom' of the possessor relation: "John's books" could be the books he owns, wrote, borrowed, recently read, is standing on to reach the top shelf: **implicit relation parameter**.

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is probably false. Now, most people in the world are too young to have grandchildren, but that doesn't make the sentence true. Grandchild-less people are simply *irrelevant* to the truth value of (17); only people in the domain of the grandparent relation matter.

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- Here  $R_a = \{b : R(a, b)\}$ , and  $dom_A(R) = \{a : A \cap R_a \neq \emptyset\}$ .

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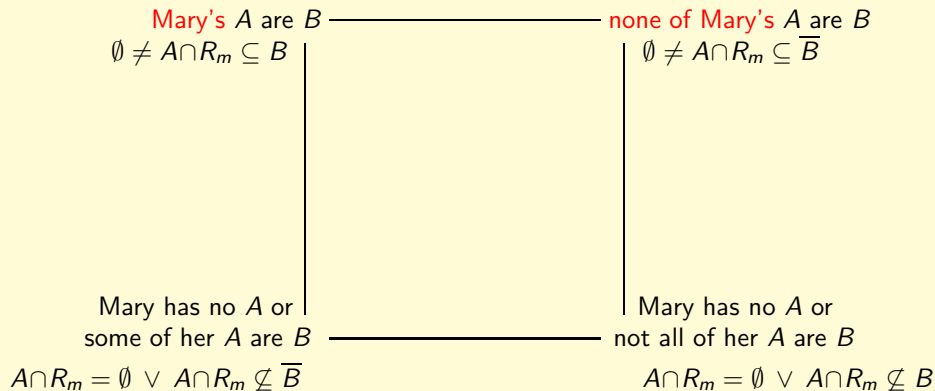
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- But this does not seem to hold for the class of quantifiers denoted by possessive determiners.
- For example, “Mary’s cars are nice” (universal reading) is  $\emptyset \neq car \cap own_m \subseteq nice$   
(so *Mary’s* has existential import), but the outer negation, “Either Mary has no cars or some car of hers is not nice”, seems hard to get with a possessive Det.

# Possessive squares 1

$Q = (\text{all of}) \text{ Mary's} = \text{Poss}(\text{all}_{\text{ei}}, \{m\}, \text{all}, R)$ .  $Q$  is  $\text{MON}\uparrow$  and weakly  $\downarrow\text{MON}$ , so it belongs in the **A** corner:



This square happens to be classical as well.

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- However: (a) the **I** and **O** corners are unnatural, and (b) all quantifiers of the form  $Q_2$  of *Mary's* have 'possessive existential import', in the sense that their truth conditions entail  $A \cap R_m \neq \emptyset$ .



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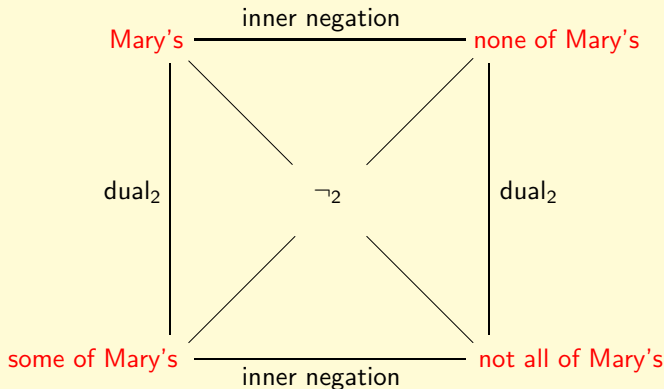
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- Now, we obtain:

# Mary's once again



This can be seen as *either*  $\text{square}(\text{Mary's})$  under the presupposition  $A \cap R_m \neq \emptyset$ , *or* as  $\text{square}^{\text{poss}}(\text{Mary's})$ , where, for possessive  $Q$ ,

$$\text{square}^{\text{poss}}(Q) = \{Q, Q\neg, \neg_2 Q, Q^{\text{d}_2}\}.$$

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- **FACT 12:** If  $Q'$  belongs to  $\text{square}^{\text{poss}}(Q)$ , then  $\text{square}^{\text{poss}}(Q') = \text{square}^{\text{poss}}(Q)$ .

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- (Of course it cannot mean that some of her friends are nice either ( $Q^{d_2}$ ), but that holds for all squares: dual is not a negation, but a form of *double* negation.)

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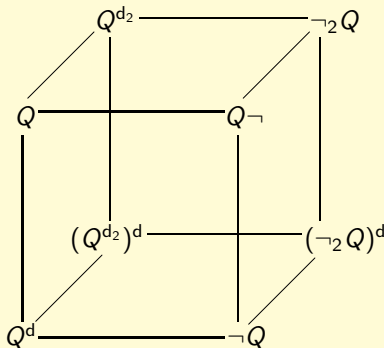
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- **FACT 15:** Of the 16 combinations of  $Q = Q_1, Q_2$  in  $Poss(Q_1, C, Q_2, R)$  ( $Q, Q \neg, \neg Q, Q^d$ ), only 8 yield distinct results.

# A cube of opposition

This allows arranging all the 'possessive oppositions' in a **cube of opposition**, with inner negation along 4 sides, dual along 4 sides, and dual<sub>2</sub> along the remaining 4:



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- Let us look at the possible squares for these.

## *square(no car's)* (existential)

$Q(A, B) \Leftrightarrow \text{Poss}(\text{no}, C, \text{some}, R)(A, B) \Leftrightarrow C \subseteq \{a: A \cap R_a \subseteq \overline{B}\}$ . By a theorem in P & W 2006 (or directly),  $Q$  is  $\downarrow\text{MON}\downarrow$ ; hence in the **E** corner. In contrast with *Mary's*,  $Q$  has no 'possessive existential import'.

each car's  $A$  were (all)  $B$

$$C \subseteq \{a: A \cap R_a \subseteq B\}$$

(cf. Fact 13)

no car's  $A$  were  $B$

$$C \subseteq \{a: A \cap R_a \subseteq \overline{B}\}$$

at least one of some car's  $A$  was  $B$

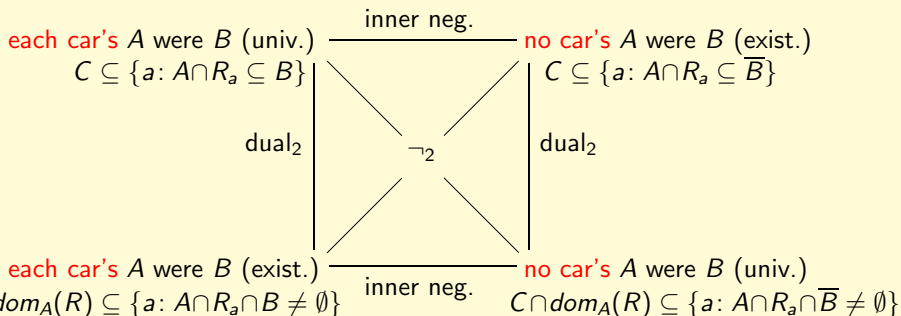
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not all of some car's  $A$  were  $B$

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# $square^{poss}(no\ car's)$ (existential)

The possessive square for *no car's* (existential) differs at the **I** and **O** corner. Interestingly, at the **O** corner we get the universal reading of *no car's*; the one I claimed was sometimes reasonable for *No student's books were returned*. Analogously for the **A** and **I** corners:



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- A theorem in P & W ch. 7 describes how the monotonicity properties of  $Q_1$  and  $Q_2$  determine those of  $Poss(Q_1, C, Q_2, R)$ .
- From it, you might easily think that if  $Q_1$  and  $Q_2$  are both doubly monotone, so is  $Poss(Q_1, C, Q_2, R)$ . But although this is usually true, there are exceptions:
- **FACT 16:**  $Poss(no, C, all, R)$ , e.g., the universal reading of *no students'*, is  $\text{MON}\downarrow$  but neither  $\uparrow\text{MON}$  nor weakly  $\downarrow\text{MON}$ .
- But we saw that, in spite of this, all eight quantifiers occupy a unique corner in *cube(no students')* (since sufficiently many of them are doubly monotone).